

Universidad Politécnica de Madrid ETS de Ingenieros Informáticos



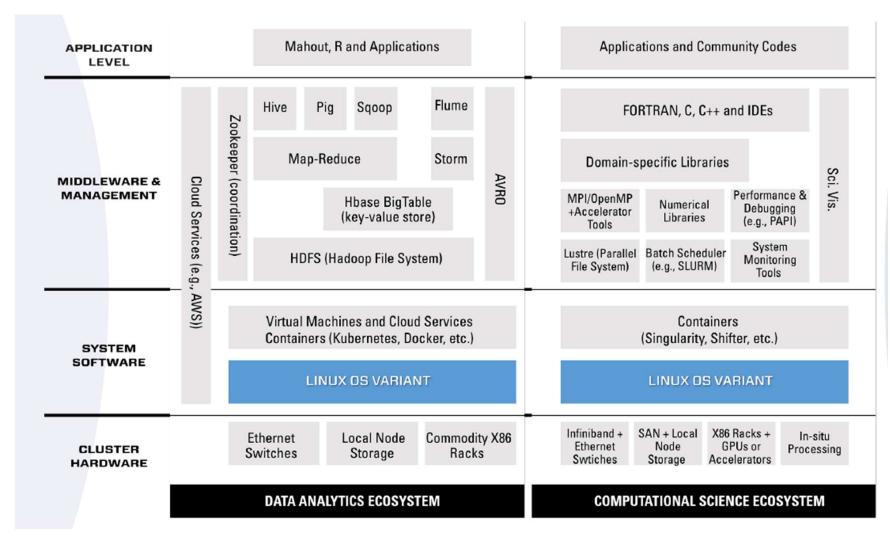


María S. Pérez mperez@fi.upm.es

- Introduction
- Context
- General convergence HPC-Big Data problem
- Convergence HPC-Big Data at storage level
- Study and evaluation
- Conclusions

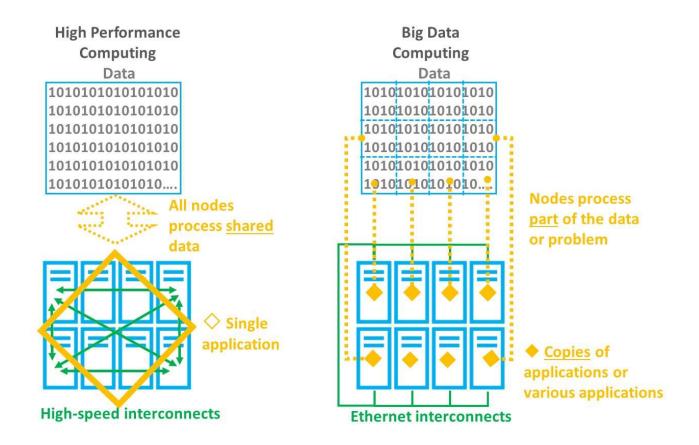


Divergence DA-CS



Source: Big Data and Extreme-Scale Computing, BDEC

Divergence at architecture level



Source: "Creating synergies across HPC & Big Data platforms", BDVA-ETP4HPC White Paper

Historical differences BD-HPC

	Typical workload	Design principles
Big Data	Data-intensive applications Most time is used for I/O and data management	Optimized for cost Less priority for performance
HPC	Compute-intensive applications Most time is used for computing	Optimized for performance Less priority for cost



European Commission > Strategy > Digital Single Market > Policies >

Digital Single Market

POLICY

The European High-Performance Computing Joint Undertaking - EuroHPC

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Useful links

The European High-Perfomance Computing Joint
Undertaking (EuroHPC JU) will pool European resources to
develop top-of-the-range exascale supercomputers for
processing big data, based on competitive European
technology.

About HPC

Policies +

Blog posts

News

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Sweden

Norway

Finland

Norway

Finland

Norway

Finland

Norway

Finland

Norway

Finland

Dementer

Lithuania

Lithuania

Mincoara

Mincoara

Mincoara

Lithuania

Norway

France

Saline See

Estonia

Lithuania

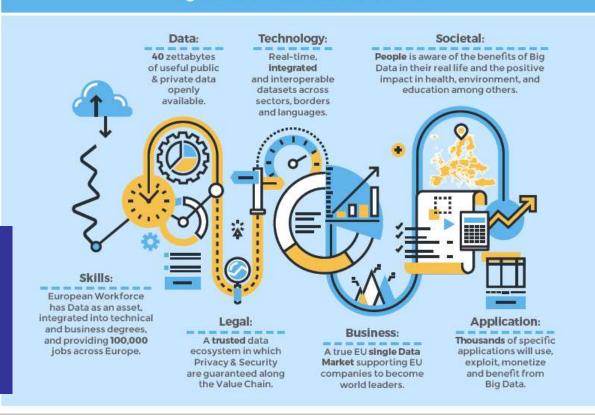
Mincoara

Mi

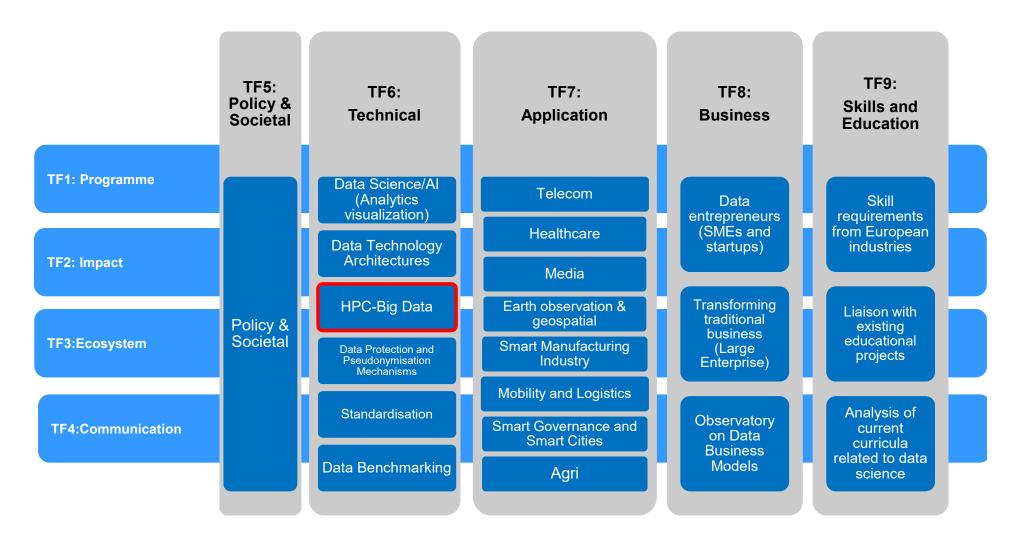
BDVA (~200) members include large industries, SMEs, research organisations and data users and providers to support the development and deployment of the EU Big Data Value Public-Private Partnership with the European Commission

BDVA focuses its activities on updating the multi-annual roadmap and on providing regular advice to enable the European Commission to prepare, draft and adopt the periodic Work Programmes, as well as on delivering Data Innovation Recommendations, developing Big Data Value Ecosystem, guiding Standards, and, facilitating Know-how exchange.

Big Data Value Vision for 2020



BDVA Task Forces



EuroHPC Joint Undertaking

https://eurohpc-ju.europa.eu

EuroHPC - Activities



Infrastructure & Operations	R&I, Applications & Skills	JU Admin/Running costs
~270	min 180	10
200	196	10

Infrastructure + Operations

Acquisition of 2 pre-exascale machines and several (tbd) mid-range machines

Applications & Skills + R&I

R&I, exascale technologies and systems (incl. low-power processor); applications

JU Admin/running costs

JU Operation: 2019 to 2026

486 m€	10	min 180	~270	
486 m€	10	~186	~290	
972 m€	20	392	560	
422 m€	2	~420 (in kind)	0	

EC	
Participating States	
Total	
Private Members	

https://eurohpc-ju.europa.eu/documents/EuroHPC-Work-Plan-2019.pdf

CABAHLA-CM: Convergence BD-HPC: From sensors to applications

- Project funded by the Community of Madrid, grants for R&D activities between research groups in technology and biomedicine (2019-2022)
- 4 groups:
 - ArTeCS, Universidad Complutense de Madrid
 - ARCOS, Universidad Carlos III de Madrid
 - SciTrack, Centro de Investigaciones Energéticas, Medioambientales y Tecnológicas
 - OEG, Universidad Politécnica de Madrid
- Goal: Improve the integration of HPC and Big Data paradigms
 - Computing and data intensive platform
 - Two use cases: capturing and modeling sensor data for the prediction of solar radiation with high spatio temporal resolution and processing massive data in brain's medical images

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Data Intensive

Lower <---->higher

Traditional Big Data



Data-intensive workloads

[Example] Inferring new insights from big data-sets e.g. pattern recognition across suppliers, consumers, eth for data-driven insights and innusation

Extreme Data Analytics



Compute- and Data intensive workloads: [Example] Reshaping healthcare through advanced analytics and artificial intelligence — leading to predictive and personalized medicine

Enterprise IT



'Regular' workloads

[Example] Running the enterprise – HR, Legal, Payroll, finance, etc.

HPC



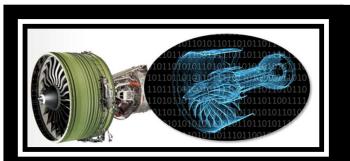
Compute-intensive workloads

[Example] Modelling and simulating focusing on interaction amongst parts of a system and the system as a whole e.g. product design

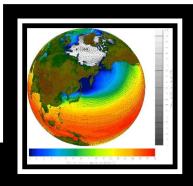
Compute Intensive Lower <---->higher

Source: Subgroup HPC-BD BDVA

Applications/Use cases

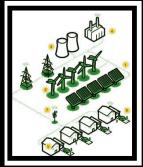


















Source: Subgroup HPC-BD BDVA

HPC, Big Data and Deep Learning

	Supercomputing (SC)	Deep Learning (DL)	Big Data (BD)	
Boundary Interaction Services	In-house, commercial & OSS applications [e.g. Paraview], Remote desktop [e.g. Virtual Network Computing (VNC)], Secure Sockets Layer [e.g. SSL certificates]	Framework-dependent applications [e.g. NLP, voice, image], Web mechanisms [e.g. Google & Amazon Web Services], Secure Sockets Layer [e.g. SSL certificates]	Framework-dependent applications [e.g. 2/3/4-D], Secure Sockets Layer [e.g. SSL certificates]	
Processing Services	Domain specific frameworks [e.g. PETSc], Batch processing of large tightly coordinated parallel jobs [100s – 10000s of processes communicating frequently with each other]	DNN training & inference frameworks [e.g. Caffe, Tensorflow, Theano, Torch], DNN numerical libraries [e.g. dense LA], DNNs, statistics, diagnostics [e.g. ?]	Machine Learning (traditional) [e.g. Mahout, Scikit-learn], Analytics / Statistics [e.g. Python, ROOT, R, Matlab, SAS, SPSS, Sci-Py], Iterative [e.g. Apache Hama], Interactive [e.g. Dremel, Drill, Tez, Impala, Shark, Presto, BlinkDB], Batch / Map-Reduce [e.g. MapReduce, YARN, Sqoop, Spark], Real-time/streaming [e.g. Flint, YARN, Druid, Pinot, Storm, Samza, Spark streaming]	
Model / Information Managemer Services	Data Storage: Parallel File Systems [e.g. Lustre, GPFS, BeeGFS, PanFS, PVFS], I/O libraries [e.g. HDF5, PnetCDF, ADIOS]	Data Storage [e.g. HDFS, Hbase, Amazon S3, GlusterFS, Cassandra, MongoBD, Hana, Vora]	Serialization [e.g. Avro], Meta Data [e.g. HCatalog], Data Ingestion 2. Integration [e.g. Flume, Squop, Apacho Nifi, Elastic Logstash, Kafka, Talend, Pentaho], Data Storage [e.g. HDFS, Hbase, Amazon S3, GlusterFS, Cassandra, MongoBD, Hana, Voral. Cluster Memt [e.g. YARN, MFSO]	
Communica Services	on Messaging & Coordination [e.g. MPI/PGAS, direct fabric access], Threading [e.g. OpenMP, task-based models]	Messaging & Coordination [e.g. Machine Learning Scaling Library (MLSL)]	Messaging [e.g. Apache Kafka (streaming)]	
Workflow / Task Service	Conventional compiled languages [e.g. C/C++/Fortran], Scripting languages [e.g. Python, Julia,]	Scripting languages [e.g. Python]	Workflow & Scheduling [e.g. Oozie], Scripting languages [e.g. Keras, Mocha, Pig, Hive, JAQL, Python, Java, Scala]	
System Managemer & Security Services	Domair numerical libraries [e.g. PETSc, ScaLAPACK, BLAS, FFTW.]. Performance & debugging [e.g. DDT, Vtune, Vampir], Accelerator APIs [e.g. CUDA, OpenCL, OpenACC] Data Protection [e.g. System AAA, OS/PFS file access control] Batch scheduling [e.g. SLURM],	Batching for training [built into DL frameworks], Reduced precision [e.g. Inference engines], Load distribution layer [e.g. Round robin/load balancing for inference []. Accelerator APIs [e.g. CUDA. OpenCI]. LA numerical libraries [e.g. BLAS, LAPACK,]	Distributed Coordination [e.g. ZooKeeper, Chubby, Paxos], Provisioning, Managing & Monitoring [e.g. Ambari, Whirr, BigTop, Chukwa], SVM systems [e.g. Google Sofia, libSVM, svm py,], Hardware Optimization Libraries [e.g. cuDNN, MKL-DNN, etc.)]	
Services Cluster management [e.g. OpenHPC], Container Virtualization [e.g. Docker], Operating System [e.g. Linux OS Variant]		Virtualisation [e.g. Dockers, Kubernetes, VMware, Xen, KVM, HyperX], Operating System [e.g. Linux (Red Hat, Obuntu, etc.), (Windows?)]	Virtualisation [e.g. Dockers, Kubernetes, VMware, Xen, KVM, HyperX], Operating System [e.g. Linux (кеанат, Obuntu, etc.), windows	
Infrastruct	Local storage Servers [e.g. CPU & Network [e.g.	Local storage [e.g. Local storage or NAS/SAN] Servers [e.g. CPU & Memory) [Gen Purpose CPU + GPU/FPGA, TPU] Network [e.g. Ethernet‡]	Local storage [e.g. Direct attached Storage] Servers [e.g. CPU & Memory, [Gen Purpose CPU hyper-convergent nodes] Network [e.g. Ethernet fabrics]	

Source: "Creating synergies across HPC & Big Data platforms", BDVA-ETP4HPC White Paper

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HPC at storage level

POSIX file system - Random reads and writes to file - Folder/file hierarchies - Permissions - Atomic file rename - Multi-user protection

POSIX

Random reads and writes to file Folder/file hierarchies
Permissions
Atomic file rename
Multi-user protection



POSIX

Random reads and writes to file



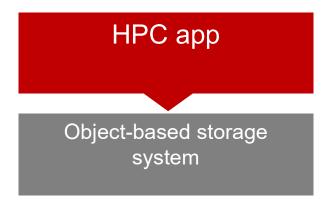
POSIX

Random reads and writes to file Objects

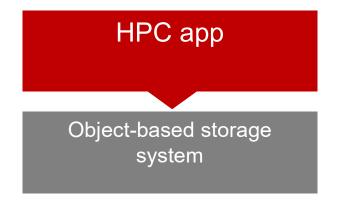
HPC at storage level

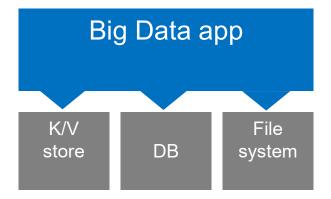
HPC app

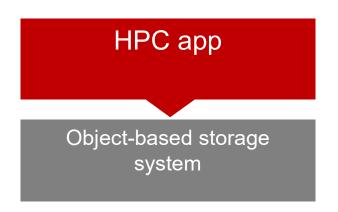
Object-based storage system

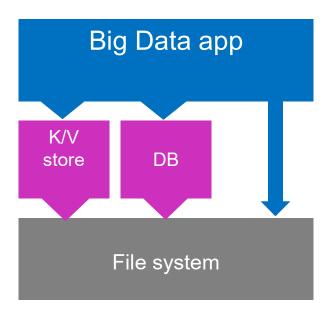


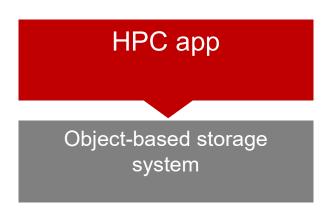
Big Data app

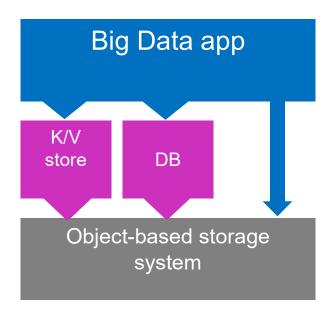


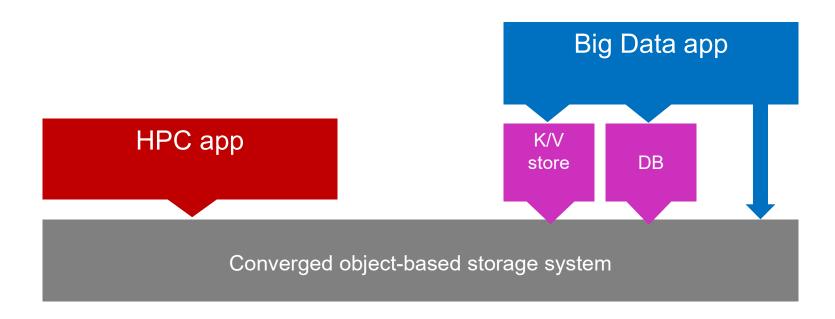




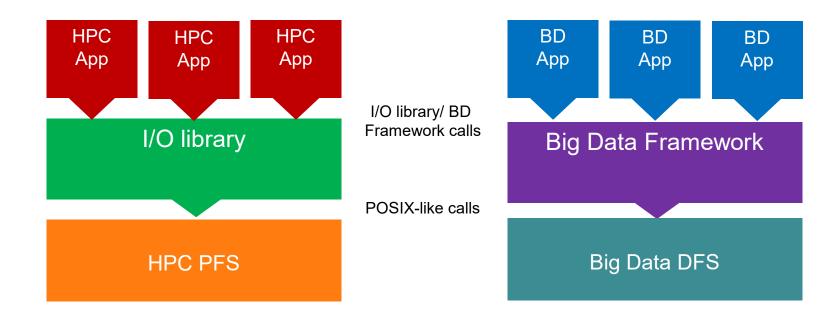




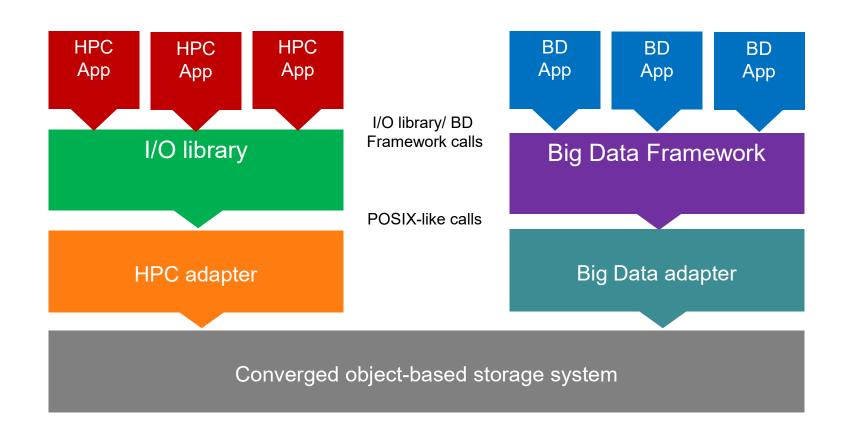




Actual storage stack



Actual storage stack



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Object-oriented primitives

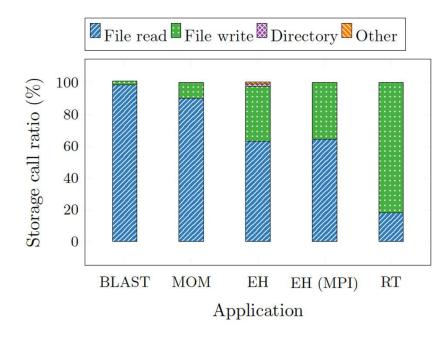
- Object access: random object read, object size
- Object manipulation: random object write, truncate
- Object administration: create object, delete object
- These operations are similar to those permitted by the POSIX-IO API on a single file
- Directory-level operations do not have their object-based storage counterpart (flat nature of these kinds of systems):
 - I ow number of them
 - Emulated using the scan operation (far from optimized, but compensated by the gains permitted by using a flat namespace and simpler semantics)

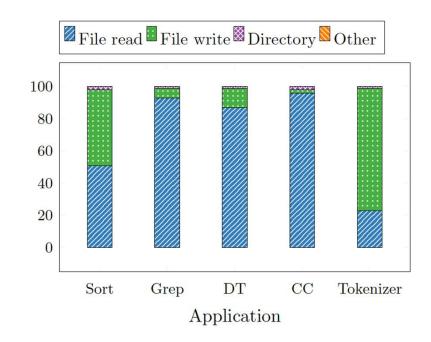
Representative set of HPC/BD Apps

Plataform	Application	Usage	Total reads	Total writes	R/W ratio	Profile
	mpiBLAST	Protein docking	27.7 GB	12.8 MB	2.2*10^3	Read-intensive
HPC/MPI	MOM	Oceanic model	19.5 GB	3.2 GB	6.09	Read-intensive
	ECOHAM	Sediment propagation	67.4 GB	71.2 GB	0.94	Balanced
	Ray Tracing	Video processing	0.4 GB	9.7 GB	4.1*10^-2	Write-intensive
	Sort	Text processing	5.8 GB	5.8 GB	1.00	Balanced
Cloud/Spark	Connected Component	Graph processing	13.1 GB	71.2 MB	1.8*10^2	Read-intensive
	Grep	Text processing	55.8 GB	863.8 MB	66.14	Read-intensive
	Decision Tree	Machine Learning	59.1 GB	4.7 GB	12.57	Read-intensive
	Tokenizer	Text processing	55.8 GB	235.7 GB	0.23	Write-intensive

Pierre Matri, Yevhen Alforov, Álvaro Brandón, María S. Pérez et al. Mission possible: Unify HPC and Big Data stacks towards application-defined blobs at the storage layer. Future Generation Computer Systems, In press.

Ops distribution





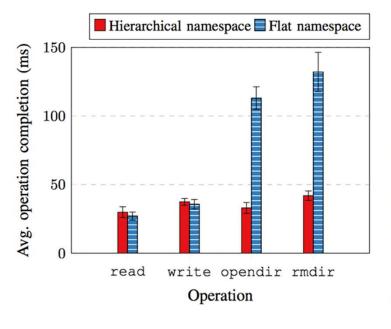
Directory operations (BD App)

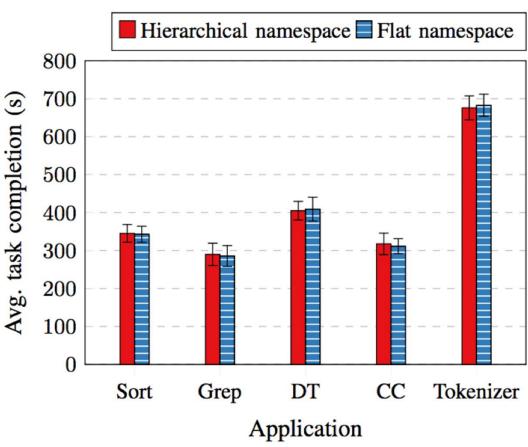
Operation	Action	Oper. count
mkdir	Create directory	43
rmdir	Remove directory	43
opendir (Input data directory)	Open/List directory	5
opendir (other directories)	Open/List directory	0

Original operation	Rewritten operation
create(/foo/bar)	create(/foobar)
open(/foo/bar)	open(/foobar)
read(fd)	read(bd)
write(fd)	write(bd)
mkdir(/foo)	
opendir(/foo)	scan(/), return all files matching foo*
rmdir(/foo)	scan(/), remove all files matching foo*



Influence of directory operations

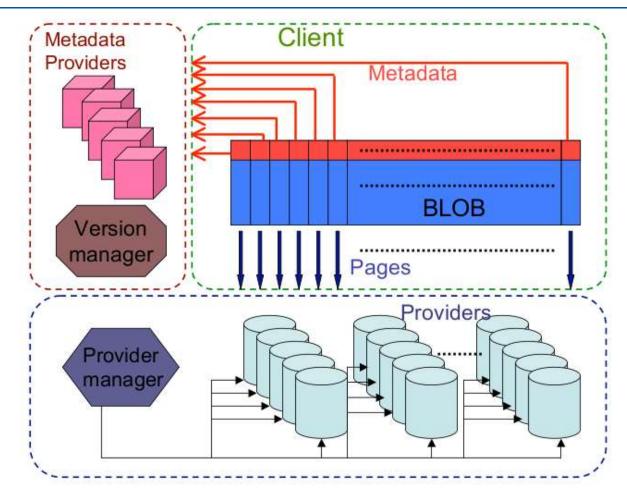




BlobSeer/RADOS vs Lustre (HPC) and HDFS/Ceph (BD)

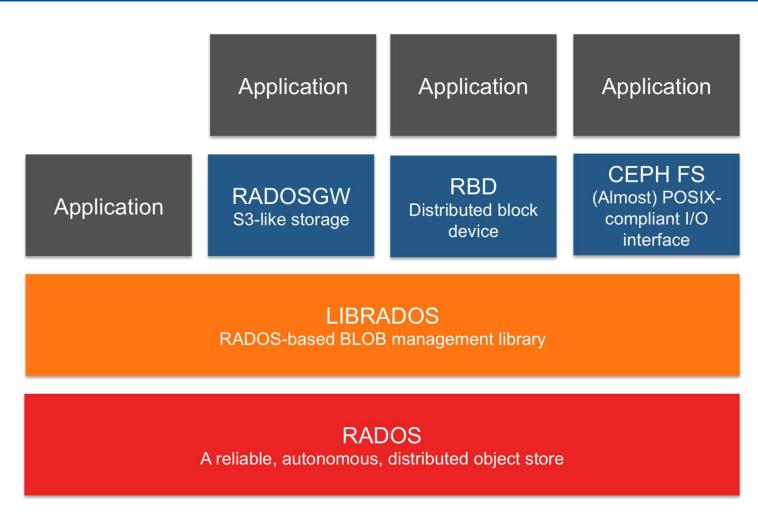
- Grid'5000 experimental testbed distributed over 11 sites in France and Luxembourg (parapluie cluster, Rennes)
- Each node: 2 x 12-core 1.7 Ghz 6164 HE, 48 GB of RAM and 250 GB HDD.
- HPC Apps: Lustre 2.9.0 and MPICH 3.2 [67], on a 32-node cluster (InfiniBand)
- BD Apps: Spark 2.1.0, Hadoop / HDFS 2.7.3 and Ceph Kraken, on a 32-node cluster (Gigabit Ethernet)

BlobSeer



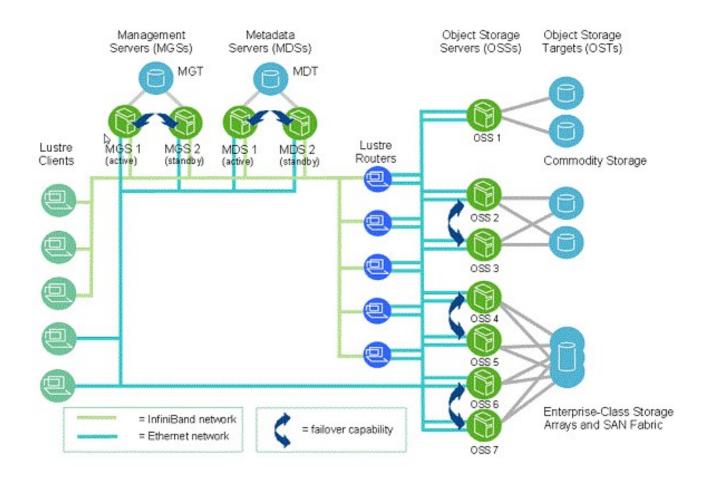
Bogdan Nicolae; Gabriel Antoniu; Luc Bougé; Diana Moise; Alexandra Carpen-Amarie. 2011. BlobSeer: Next-generation data management for large scale infrastructures. J. Parallel Distrib. Comput. 71, 2 (February 2011), 169-184.

RADOS/Ceph



ceph.com

Lustre

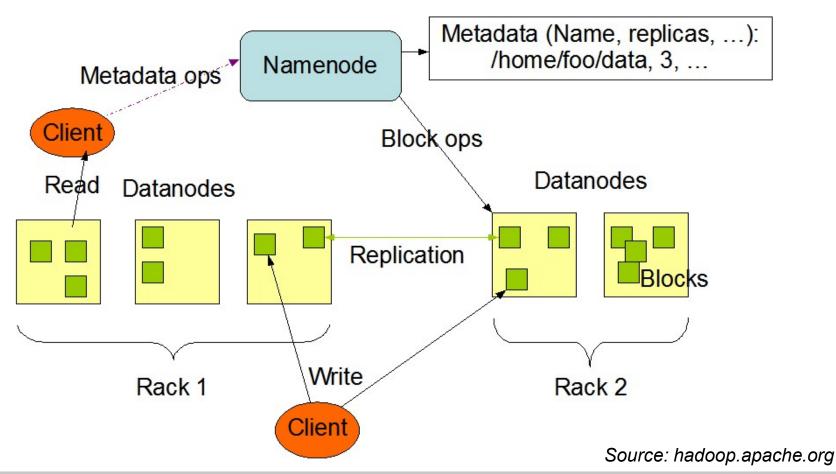


Source: lustre.org

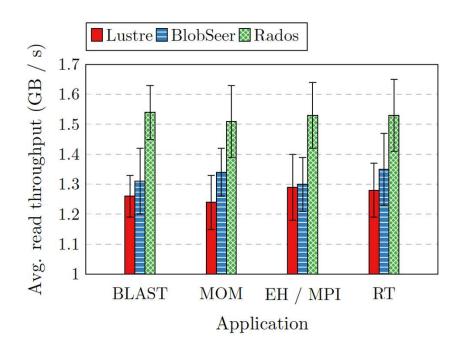


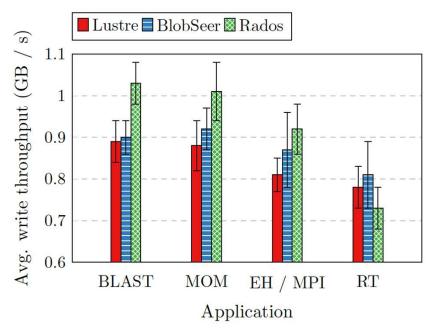
HDFS

HDFS Architecture

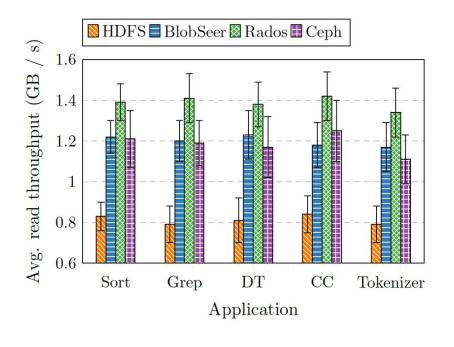


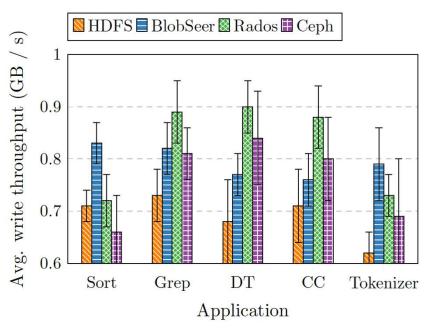
HPC Apps





BD Apps





Study insights

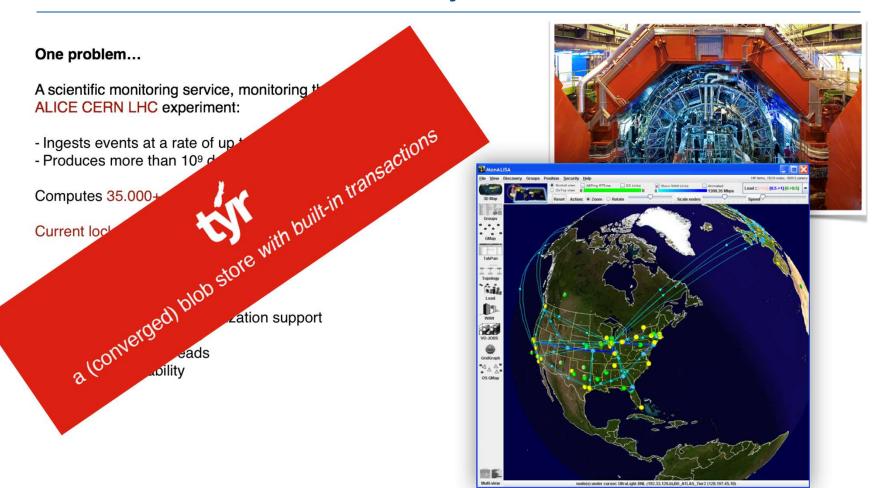
- The convergence at storage level is possible by using objectbased storage systems, achieving an improvement in the performance for both platforms (HPC and Cloud)
- By using objects, it is possible to achieve a maximum improvement of 32%
 - Mainly because of the flat namespace
 - Rados: direct reads and simple and decentralized metadata management (high performance for read operations)
 - BlobSeer: multi-version concurrency control supports high-performance write operations with highly concurrent workloads (high performance for write operations)
- Issues of both systems:
 - Although the Rados performance is excellent when the write contention is low, its lock-based concurrency control limits the performance of highly concurrent use cases.
 - The multi-version concurrency control in BlobSeer provides good performance at write level, but the **distributed metadata in BlobSeer** provides a significant read latency

Týr

- Can we achieve both systems benefits?
- Apart from combining the best features of Rados and BlobSeer, there
 is a significant set of use cases with stricter consistency semantics
 - Indexing and data aggregation (E.g., ALICE CERN LHC experiment)
 - Distributed shared logs (E.g., Computational steering + in-situ visualization)



Týr



Pierre Matri; Alexandru Costan; Gabriel Antoniu; Jesús Montes; María S. Pérez. "Týr: Blob StorageSystems Meet Built-In Transactions". SC '16 Proceedings of the International Conference for High Performance Computing, Networking, Storage and Analysis. Article n. 49, Best student paper award finalist

Týr

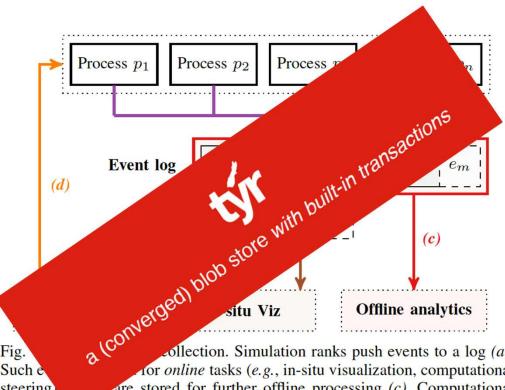


Fig. collection. Simulation ranks push events to a $\log (a)$. Such e for *online* tasks (e.g., in-situ visualization, computational steering) are stored for further offline processing (c). Computational steering in arn controls the simulation (d).

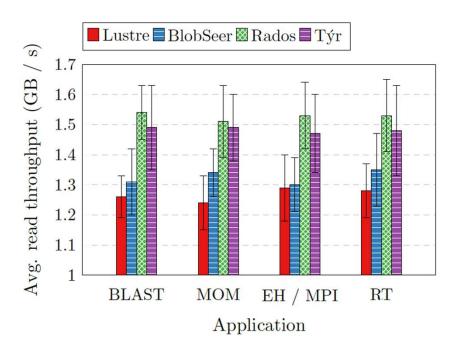
Pierre Matri; Philip Carns; Robert Ross; Alexandru Costan; María S. Pérez; Gabriel Antoniu;. "SLoG: A large-scale Logging Middleware for HPC and Big Data convergence". ICDCS'2018. pp. 1507-1512, 2018.

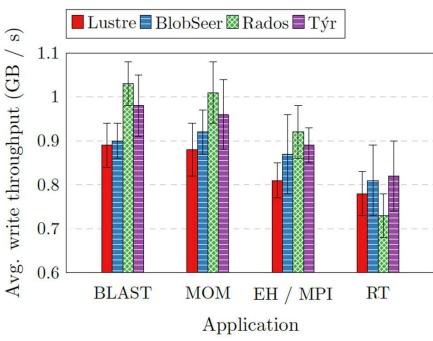
Týr design

- Predectible data distribution
 - Combining data striping and consistent hashing techniques
 - Avoiding the use of a centralized metadata server
- Transparent multi-version concurrency control
 - Versioning at chunk level
- Lightweigh ACID transactions
 - By using the transactional protocol Warp*
 - By using chains of the servers involved in transactions and dependency graphs
- Atomic transformation operations
 - Efficient read-modify-write ops
 - Particularly interesting for simple transformation operations (arithmetic, bit-wise)
 - The client does not send the new data to be written, but the modification, avoiding two-round trips
- Software prototype with around 25,000 Rust and GNU C code lines

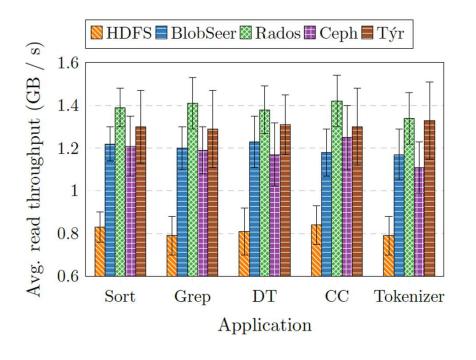
^{*} R. Escriva, B. Wong and E. Sirer. Warp: Lightweight multi-key transactions for key-value stores. arXiv preprint arXiv:1509.07815, 2015.

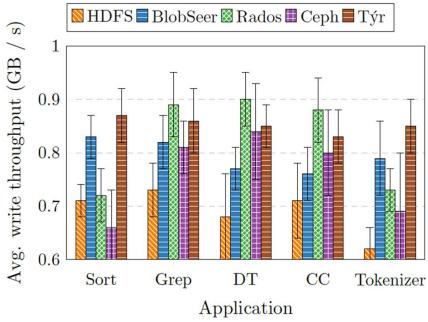
HPC Apps





BD Apps





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Conclusions

- Týr outperforms BlobSeer and traditional file systems, for both HPC and BDA applications
 - Non blocking writes, by means of multiversion concurrency control
 - Direct writes by using consistent hashing
- Týr has some limitations in comparison with Rados
 - Except for write-intensive applications, due to the efficiency of the multiversion concurrency control
 - As a result of stronger consistency guarantees (transactions)
- This is a first step for the convergence between HPC and BDA (at storage level)



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